

# Glimmers of Autonomy: Structure-Aware Reachability Analysis and Control Synthesis for Unknown Nonlinear Systems

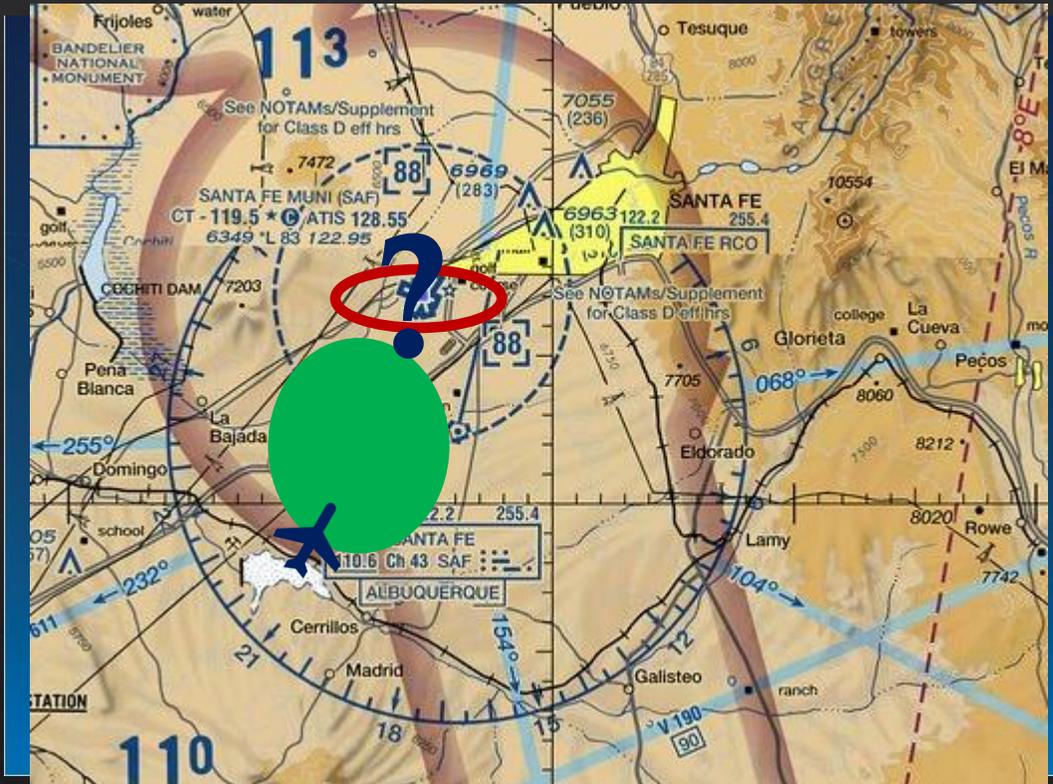
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# Maximal Understanding in the Face of Minimal Knowledge



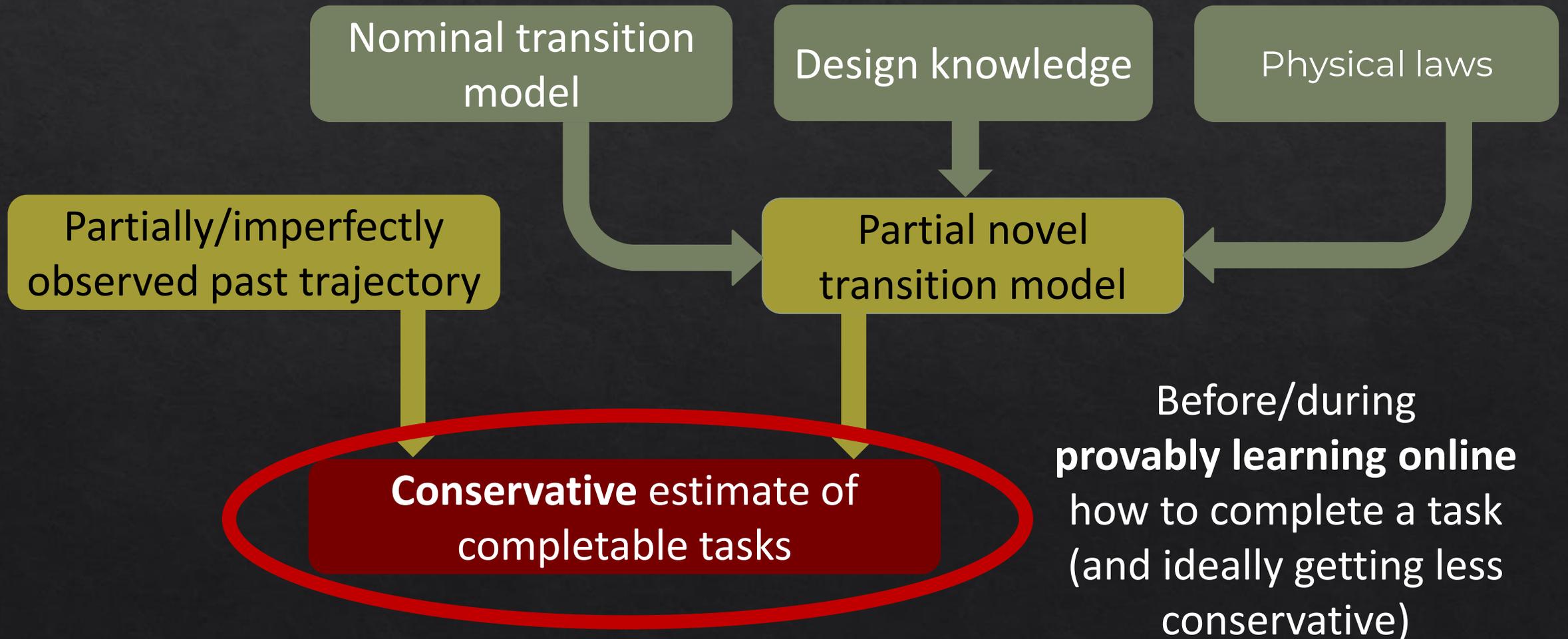
**Adaptive / Robust Control**

Reach the objective under (some) lack of knowledge of dynamics

**Objective might not be reachable**

First decide on a reachable target,  
*then* plan how to get there

# Certi fiable Capabilities



# Disasters

$$\dot{x} = f(x, u), \quad u \in U$$

Change in dynamics  
(e.g., physical damage)

Partial loss of control  
(e.g., adversarial  
takeover)

Actuator  
degradation

$$\dot{x} = \hat{f}(x, u, v), \quad (u, v) \in \hat{U}$$

# Loss of Control Authority

$$\dot{x} = f(x, u, v), \quad (u, v) \in U$$

Uncontrolled system input  
(not a disturbance)

**Two-player game:**

P1 (“controller”) wishes to reach a state

P2 (“environment”/“adversary”) wishes to obstruct P1

**Can P1 win?** (Can P1 win for any state? **Can P1 win if there is a time limit?**)

“resilience”

# Quantitative Resilience

Intuitively, a system is resilient if a player can counteract the adversary and still have some meaningful control authority

If the target state is in the guaranteed reachable set of the initial state (*guaranteed* with respect to all possible adversarial control inputs):

it also matters **how hard it is** to reach the state compared to the system with nominal dynamics.

The adversary chooses an input such that whatever the controller chooses, performance will be bad

$$r_q = \frac{\inf_{\bar{u}} T_{\bar{R}}^{\bar{u}}(x_0, x_{goal})}{\sup_u \inf_v T_{\mathcal{P}}^{(u,v)}(x_0, x_{goal})}$$

# Reach-Time Resilience

**First idea:** measure of performance – reach time

$$r_q = \frac{\inf_{\bar{u}} T_R^{\bar{u}}(x_0, x_{goal})}{\sup_v \inf_u T_R^{(u,v)}(x_0, x_{goal})}$$

**Resilience quotient = 0:** no resilience; **resilience quotient = 1:** perfect resilience

Determining **minimal time** to reach a target even for **nominal linear** dynamics with input constraints is not simple (*see last year*)

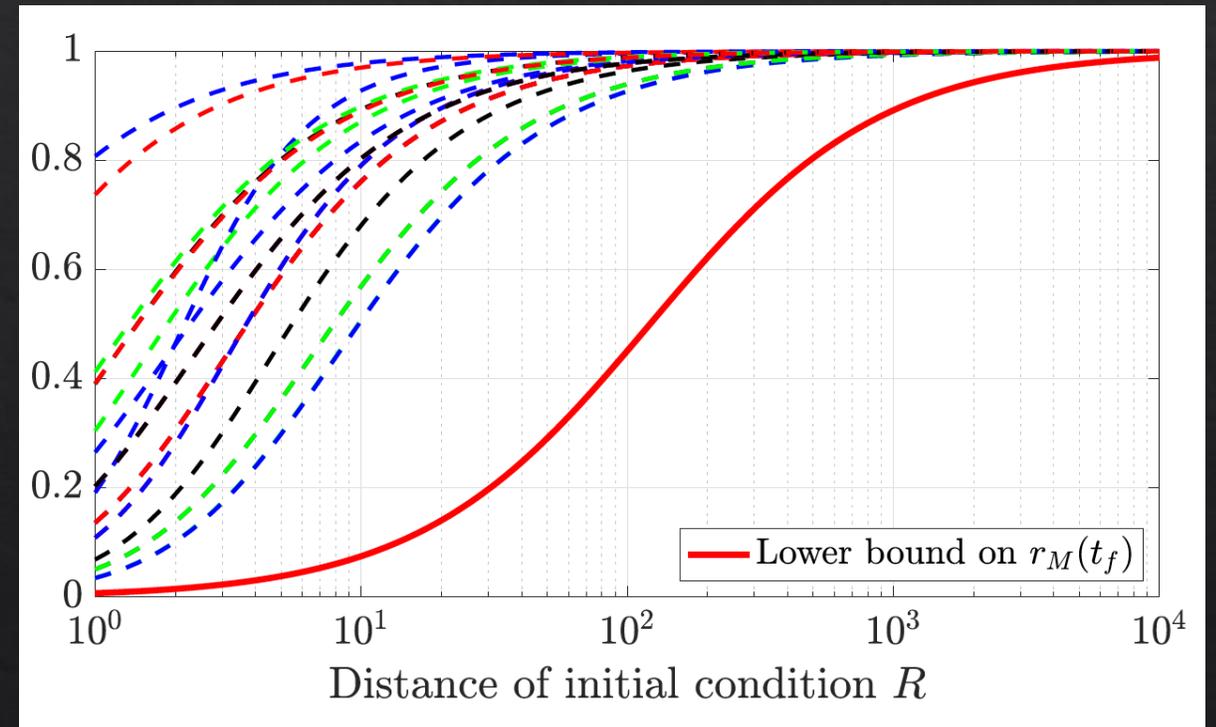
# Energy Resilience

Determining **minimal energy** to reach a target (soft input constraints) for nominal dynamics **is simple**: controllability Gramian

*for linear systems*

**Step 0:** Determine worst-case minimal energy for reachability at a particular time for *disturbed dynamics*

Slight problem: quotient of energies goes to zero the closer we start to equilibrium (player constantly fights the adversary) – difference of energies



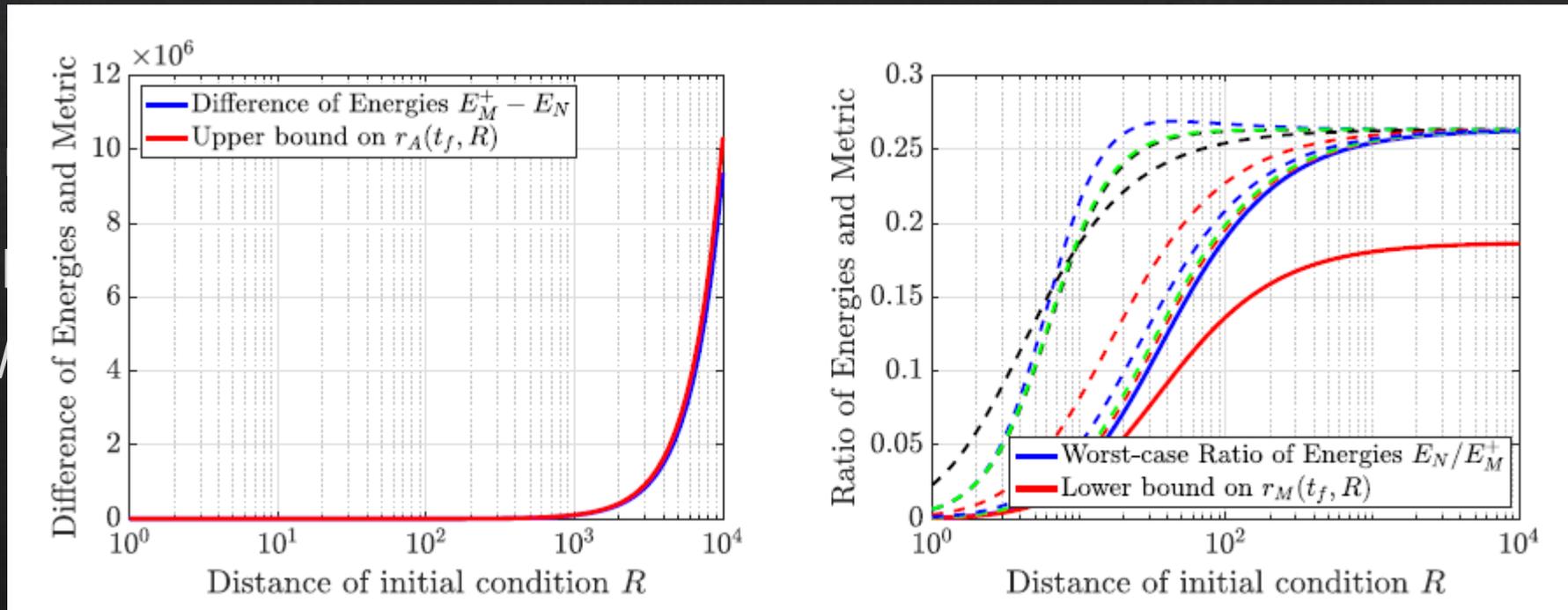
# Energy Resilience for Driftless Systems

“Disturbances” in our scenario have structure, *and* faulty/hostile actuators still expend energy

For d

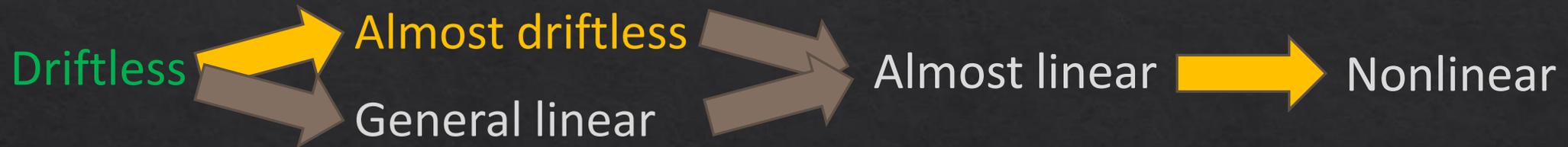
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The stronger the adversary, the larger the energy  
The stronger the controller, the smaller the energy

# Towards Nonlinearity



Almost driftless: growth or magnitude bound on a nonlinear drift term

- ◇ Idea: difficult to express optimal control input, possible to bound it  
*(first results this week)*

Nonlinear (on a compact set) = linear + bound on a nonlinear term

*Switched systems?*

# Disasters

$$\dot{x} = f(x, u), \quad u \in U$$

Change in dynamics  
(e.g., physical damage)

Partial loss of control  
(e.g., adversarial takeover)

$$\dot{x} = \hat{f}(x, u, v), \quad (u, v) \in \hat{U}$$

# Guaranteed Reachability on Manifolds

$$\dot{x} = f^?(x, u), \quad u \in U$$

$\dot{x} \in T_x \mathcal{M}$        $x \in \mathcal{M}$

Physical knowledge,  
a priori safety constraint?

After a change in dynamics, the system might not be able to reach its target using **any** control law

What is it **certifiably capable** of doing (even if we don't yet know how)?

$$R^{\mathcal{G}}(x_0) = \bigcap_{\tilde{f} \in D_{con}} R^{\tilde{f}}(x_0)$$

# Idea in Euclidean Space

- ◇ In theory, guaranteed reachability set is well-defined
- ◇ In practice, how do compute it?

**Idea:** By finding all trajectories obtained by integrating **guaranteed velocities**

$V^{\mathcal{G}}(t) = \bigcap_{\tilde{f} \in D_{con}} V^{\tilde{f}}(t)$ 
Difficult, but one level easier to underapproximate
 we will obtain at least some (if not all) guaranteed reachable states

Velocities guaranteed at time  $t$  = Online "local System ID"  
~~velocities guaranteed at start time~~  
~~"modulo" maximal system wildness~~  
"Physics and design" = Lipschitz

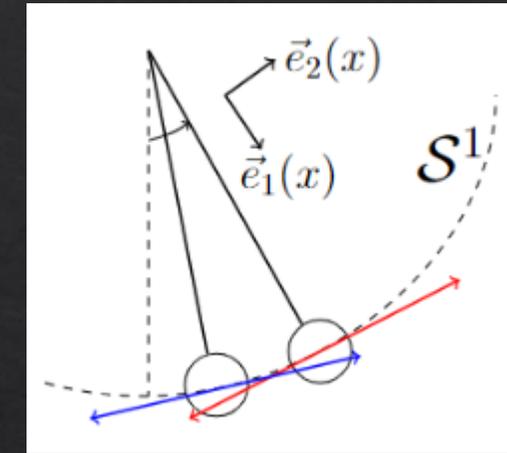
# Challenge with Manifolds

- ◇ Velocities guaranteed at different times are no longer in the same space ( $T_x M$ )
- ◇ Lipschitz constant on  $M$  is nontrivial to define

Let  $V$  be a continuous vector field on  $M$  and  $\tau$  be the parallel transport. Then  $L$  is the *classical Lipschitz constant* on  $V$  if

$$L = \sup_{\gamma} \frac{|\tau_{\gamma} V(\gamma(0)) - V(\gamma(1))|_{h_x}}{\text{Length}(\gamma)}$$

where  $\gamma : [0, 1] \rightarrow M$  varies over all  $\mathcal{C}^1$ -paths and  $\tau_{\gamma}$  is shorthand for the parallel transport along the curve  $\gamma$  from  $\gamma(0)$  to  $\gamma(1)$ .



*Why not just embed everything into Euclidean space?*

- Doesn't feel right
- Worse Lipschitz
- **No full actuation**

# Velocities on Manifolds

Idea still the same, but underapproximation of guaranteed velocities will depend on the:

Riemannian metric tensor/  
Covariant derivatives/  
Choice of connection

For flat manifolds, we recover  
previous results  
for Euclidean spaces

**Theorem 1.** Let  $f(x_0)$ ,  $G(x_0)$ ,  $L_f$ ,  $L_G$ ,  $H_x$ ,  $\Gamma_{ij}^k$ , and  $g_l^\Gamma$  for  $l \in [m]$  be defined as above. Let  $\gamma : [0, 1] \rightarrow M$  define a geodesic curve from  $x_0$  to  $x$ . Let  $\tilde{\tau}$  define the parallel transport using the flat connection. Set  $a(x) = (\|H_x^{-1}\| \|H_x\|)^{\frac{1}{2}} \|H_x\|^{\frac{1}{2}} \|[g_1^\Gamma \dots g_m^\Gamma]\|$ ,  $b(x) = (\|H_x^{-1}\| \|H_x\|)^{\frac{1}{2}} \left\| \sum_{i,j,k} \dot{\gamma}^i \Gamma_{ij}^k f^j(x_0) \vec{e}_k \right\|$ ,  $c(x) = (\|H_x^{-1}\| \|H_x\|)^{\frac{1}{2}} (L_g + \|H_x\|^{-\frac{1}{2}} L_f) d(x_0, x)$  and

$$\bar{d}(x_0, x) = \frac{\|\tilde{\tau}_{x_0}^x G^\dagger(x_0)\|^{-1} - a(x) - b(x)}{c(x)}.$$

If  $d(x_0, x) \leq \bar{d}(x_0, x)$ ,

$$\bar{\mathcal{V}}_x^G = \mathbb{B}^n(\tilde{\tau}_{x_0}^x f(x_0); \alpha(x_0, x)) \cap \text{Im}(\tilde{\tau}_{x_0}^x G(x_0))$$

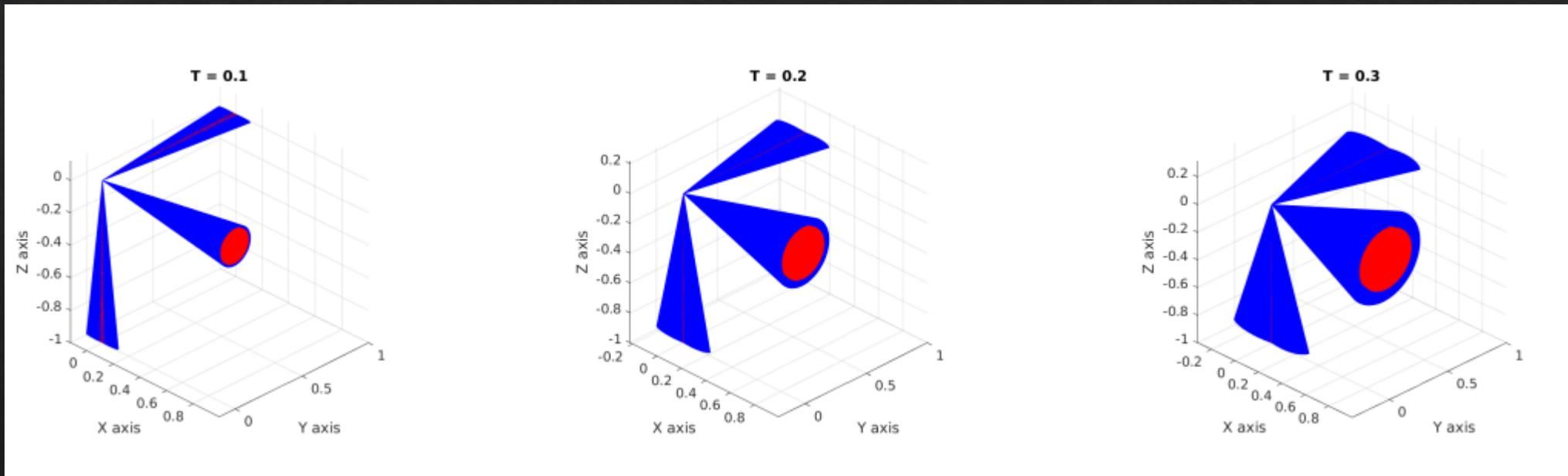
where  $\bar{\mathcal{V}}_x^G \in T_x M$ , and

$$\alpha(x_0, x) = \|\tilde{\tau}_{x_0}^x G^\dagger(x_0)\|^{-1} - (\|H_x^{-1}\| \|H_x\|)^{\frac{1}{2}} \left( \|H_x\|^{\frac{1}{2}} \|[g_1^\Gamma \dots g_m^\Gamma]\| + \left\| \sum_{i,j,k} \dot{\gamma}^i \Gamma_{ij}^k f^j(x_0) \vec{e}_k \right\| + (L_g + \|H_x\|^{-\frac{1}{2}} L_f) d(x_0, x) \right),$$

then  $\bar{\mathcal{V}}_x^G \subseteq \mathcal{V}_x^G$ .

# Numerical Example

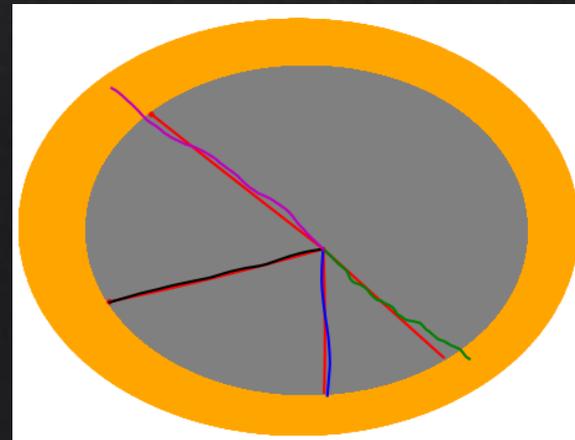
Operating on  $SO(3)$ : difficult to meaningfully think about in Euclidean space, possibly useful for applications, and difficult to even draw the guaranteed reachable set



# Learn-Control Pipeline: “Glimmers of Autonomy”

Once we have established what the system is capable of doing, we still do not know *how* to do it

- ◇ **End-to-end planning, learning and control:**
  - ◇ **Task assignment:** what task can be provably completed
  - ◇ **Real-time learning:** what do we need to know in order to be able to complete it: persistent excitation allows us to learn local dynamics
  - ◇ **Assured control:** complete the task



# Previous Work: Online “Local System ID”

When the system is at state  $x$ , applying a constant control  $u^+$  and observing the system response provides an approximation for  $f(x, u^+)$

Control-affine systems: quickly sequentially applying affinely independent controls provides an approximation for the system dynamics at  $x$

*Let  $\phi$  be the system trajectory. Let  $u^0, u^1, \dots, u^m$  be affinely independent with  $\|u^j\| \leq \delta$  for all  $j$ . Define  $x_0 = x = \phi(T)$ , and  $x_j = \phi(T + j\varepsilon)$ , where input  $u(t) = u^j$  is applied for  $t \in [T + j\varepsilon, T + (j + 1)\varepsilon)$ . Then*

$$\left\| f(x, \bar{u}) - \sum_{j=0}^m \lambda_j \frac{x_{j+1} - x_j}{\varepsilon} \right\| \leq K\varepsilon \frac{4m^{3/2} + \delta}{\delta},$$

*for all  $\bar{u}$ , where  $\bar{u} = \lambda_0 u^0 + \dots + \lambda_m u^m$ .*

# Guarantees in Online Planning

**Previously:** If we know the local dynamics, we can choose a control input that **appears to** work well right now. Goodness function encodes trajectory quality: the higher its value, the better the direction of the system.

**No guarantees! How to even choose the goodness function?**

**Step 1:** choose a target guaranteed to be reachable

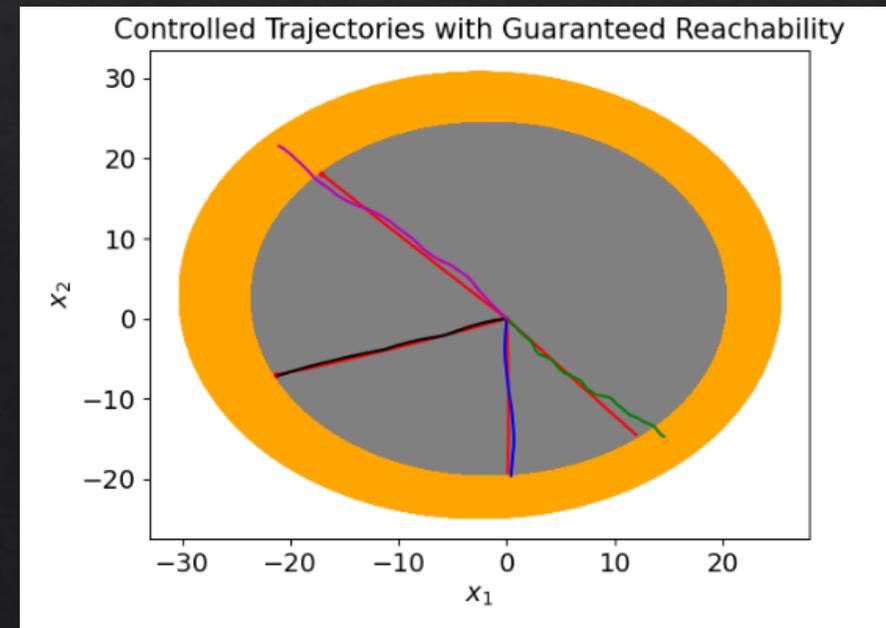
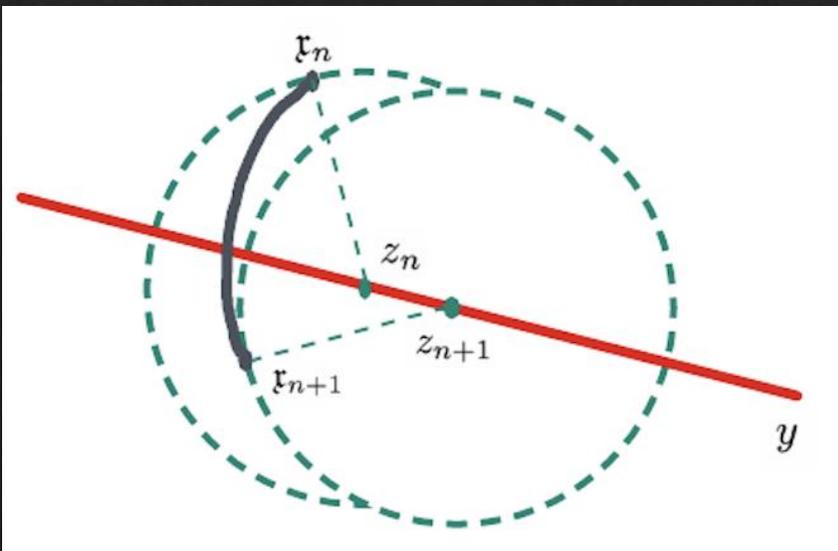
**Step 2:** guaranteed reachability computations also produce a **guaranteed trajectory**

**Step 3:** **approximately** follow the guaranteed trajectory

# Control Design

**Step 3a:** perform online learning (a previous result bounds the distance from the trajectory)

**Step 3b:** choose a point on the trajectory and use a control input that currently approximately (with known error bounds) moves the system towards that point



# Teaser: Learning from Multiple Partial Trajectories

Multiple agents operating at the same time can collect data in parallel.

**Can we obtain information about global (or “less local”) dynamics?**

Arbitrarily many agents sampling arbitrarily close: trivial

Bounded number of agents:

- ◆ **Where to place them?**
- ◆ **Interpolation with error bound quantification**

With good placement, uncertainty does not grow indefinitely – long-term guarantees?

# Teaser: Complicated Objectives

Only objective currently: **reachability/stabilization**

In progress (in some sense): **safety, trajectory tracking**

As means towards reachability, through self-designed  
**waypoints**

More complicated missions: **hybrid** (in some sense; even if the dynamics are not hybrid, the reachability specification/coordinate system might change)

**First approach:** set waypoints such that at each waypoint, the next waypoint is provably reachable (local guarantees, global heuristics)

# Medium-Term Goal (Last Year): Capabilities

“Almost driftless” with disturbance

**Combination of scenarios:** e.g., partially unknown dynamics with partial loss of control (Some work already done: actuator degradation with disturbances)

## ◇ End-to-end planning, learning and control:

- ◇ **Task assignment:** what task can be provably completed
- ◇ **Real-time learning:** what do we need to know in order to be able to complete it
- ◇ **Assured control:** complete the task

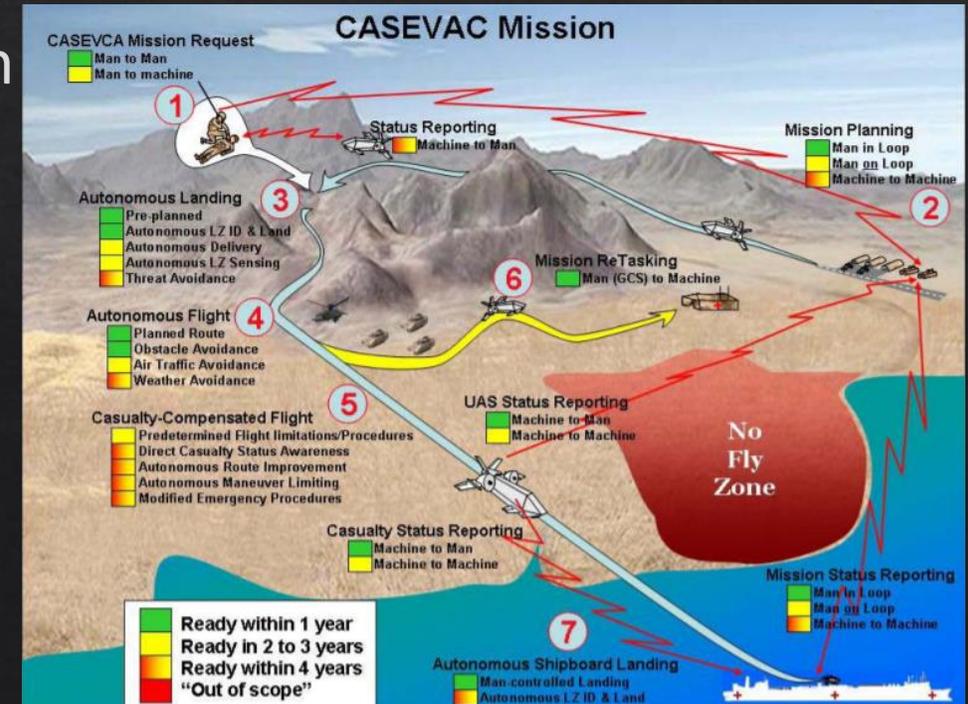
Multiple trajectories, manifolds

- ◇ (More) **physics-based and design-based results:** exploitation of significant unchanged prior knowledge for better estimation

# Long-Term Goals: Validation

- ◇ Using sensors and perception to recognize fault type and gather information
  - ◇ Fault detection; sensor fusion; state estimation
- ◇ Complex missions in high-fidelity simulation
  - ◇ Concurrent sensing and actuation faults
  - ◇ Noise+hostile action
- ◇ Onboard implementation
  - ◇ Time-delayed control
  - ◇ Real-time computation

ONR



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